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New U.S. Patent Application
Title: SPLIT-KEY KEY-AGREEMENT PROTOCOL
Inventor and Address:

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Sir:

We enclose the following papers for filing in the United States Patent and Trademark Office in connection with the above patent application.

1. A check for \$690.00 representing the filing fee.
2. Application - 8 pages, including 1 independent claim and 11 claims total.
3. Drawings - 2 sheets of informal drawings containing 2 figures.

This application is being filed under the provisions of 37 C.F.R. § 1.53(f). Applicant awaits notification from the Patent and Trademark Office of the time set for filing the Declaration.

Applicant claims the right to priority based on Canadian Patent Application No. 2,277,633, filed on July 19, 1999.

Please accord this application a serial number and filing date.

FINNEGAN, HENDERSON, FARABOW, GARRETT & DUNNER, L.L.P.

Assistant Commissioner for Patents

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The Commissioner is hereby authorized to charge any additional filing fees due and any other fees due under 37 C.F.R. § 1.16 or § 1.17 during the pendency of this application to our Deposit Account No. 06-0916.

Respectfully submitted,

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SPLIT-KEY KEY-AGREEMENT PROTOCOL

The present invention relates to the field of key agreement protocols in cryptographic systems.

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BACKGROUND OF THE INVENTION

Traditionally, entities communicated on paper and were able to ensure privacy in many ways. The transition from paper to electronic media however, has created the need for electronic privacy and authenticity. In cryptographic schemes, the entities use primitives, which are mathematical operations together with encoding and formatting techniques to provide security. For each scheme the parties participating in the scheme normally agree upon or exchange certain information before executing the scheme function. The specific information that needs to be agreed upon is detailed for each scheme. Such agreement may be achieved by any means suitable for the application. It may be implicitly built into the system or explicitly achieved by some sort of exchange of information with or without involvement from other parties. In particular, parties often need to agree on parameters and obtain each other's public keys. For proper security, a party needs to be assured of the true owners of the keys and parameters and of their validity. Generation of parameters and keys needs to be performed properly and, in some cases, verification needs to be performed.

In general, the different types of schemes may be defined as follows. Key agreement schemes, in which two parties use their public, private key pairs and possibly other information, to agree on a shared secret key. A signature scheme with appendix is a scheme in which one party signs a message using its private key and any other party can verify the signature by examining the message, the signature, and the signer's cross corresponding public key. In signature schemes with message recovery, one party signs a message using its private key and any other party can verify the signature and recover the message by examining the signature and the signer's corresponding public key. Finally, in encryption schemes, any party can encrypt a message using the recipient's public key and only the recipient can decrypt the message using its corresponding private key.

An example of a key derivation scheme is the MQV (Menezes-Qu-Vanstone). In the MQV scheme, a shared secret value is derived from one party's two key pairs and another

party's two public keys where all the keys have the same discrete log (DL) parameters. In this generalized MQV scheme, it is assumed that the shared secret value is that which is shared between two parties.

However, where each party or entity consists of a collection of parties say $A = \{A_1, A_2, \dots, A_n\}$ and $B = \{B_1, B_2, \dots, B_m\}$ where m is not necessarily equal to n and at least one of m or n is at least two (that is, not both A and B consist of one individual), it is difficult to implement the generalized MQV scheme if these two entities wish to establish a common key in order to communicate privately.

10 SUMMARY OF THE INVENTION

Accordingly, the present invention seeks to provide a solution to the problem of establishing a common key for private communication between entities wherein the entities include a collection of sub entities.

An advantage of the present invention is that all members of each entity must participate in the scheme and no subcollection of either entity can impersonate its entire entity.

In accordance with this invention there is provided a method for generating a shared secret value between entities in a data communication system, one or more of the entities having a plurality of members for participation in the communication system, each member having a long term private key and a corresponding long term public key, the method comprising the steps of:

- (a) generating an entity long term private key and corresponding entity long term public key for each entity by combining the long term private and public keys of each members of the entity.
- (b) generating a short term private and a corresponding short term public key for each of the members;
- (c) exchanging short term public keys of the members within an entity;
- (d) for each member:
 - i. computing an intra-entity shared key by mathematically combining said short term public keys of each said member;
 - ii. computing an intra-entity public key by mathematically combining its short - term private key, the long term private key and said intra-entity shared key;

- (e) for each entity combining intra-entity public keys to derive a group short-term public key;
- (f) each entity transmitting its intra-entity shared key and its group short term public key to said other entities; and
- 5 (g) each entity computing a common shared key K by combining its group short term public key, with the intra-entity shared key, and an entity long term public key received from the other entity.

10 BRIEF DESCRIPTION OF THE DRAWINGS

A preferred embodiment of the invention will now be described by way of example only with reference to the accompany drawings in which:

Figure 1 is a schematic diagram of a communication system; and

Figure 2 is a schematic diagram illustrating the steps of a protocol to establish a common key.

15 DETAILED DESCRIPTION OF THE PREFERRED EMBODIMENTS

Referring to figure 1, a schematic diagram of a communication system is shown generally by numeral 10. The system 10 includes a first entity A (12) and a second entity B (14) that exchange data over a communication channel 16. Each of the entities A and B include members

20 $A_1, A_2 \dots A_n$, and $B_1, B_2 \dots B_n$, respectively. For convenience, the embodiment described has two members A_1, A_2 and B_1, B_2 although it will be appreciated that typically each entity will have several members. It is assumed the entities A and B include processors for performing cryptographic operations and the like. The members A_1, A_2 may for example be a first group of users on a local area network (LAN) that wish to communicate securely with a second group of

25 users B_1, B_2 on a second LAN or even on the same LAN. In either case the computations may be performed for the entities A (12) and B (14) by for example a LAN server 18 or the like, provided that each member has its own secure boundary.

Each entity and its associated members A_i, B_i have been initialized with the same system parameters. The system parameters for this exemplary protocol are an elliptic curve point P,

30 which is the generating point of an elliptic curve over F_2^m of order n . Additionally, each of the members is initialized with respective long-term public and private key pairs. That is, each of

the members A_i has long term private and public key pairs (a_i, a_iP) and each of the member B_i have long term private and public key pairs (b_i, b_iP) , respectively.

Each of the entities A, B generates respective long-term public keys derived from the long-term public keys of each of its members. The long-term private key a of the entity A is then
 5 $(a_1 + a_2 + \dots a_n)$ and its corresponding long-term public key, aP , is $(a_1 + a_2 + \dots a_n) P$. In the present example the key pair (a, aP) of entity A is $(a_1 + a_2) ; (a_1 + a_2)P$. Similarly, for entity B its long-term private key b is $(b_1 + b_2)$ and its corresponding long-term public key is $bP (b_1 + b_2) P$. The entity long-term public keys aP, bP can be computed by summing the members public keys. The entity public keys are published by the respective entities, and if appropriate certified by a
 10 trusted authority or CA trusted by all of the entities.

Typically, entities A (12) and B (14) wish to agree upon a common key, which may then be used for subsequent cryptographic communications between the activities.

Referring thus to figure 2, a schematic diagram of an embodiment of a suitable protocol is shown generally by numeral 40. The member A_1 generates a random value x_1 (its short-term private key, also known as ephemeral or session key) and computes a corresponding value x_1P
 15 (its short-term public key); similarly, member A_2 generates a random value x_2 and computes a corresponding value x_2P . Preferably $0 < a_i < n-1$ and $0 < x_i < n-1$. Next, the members of the entity A exchange their session public keys x_iP . In the present example, A_2 and A_1 exchange their session public keys x_1P and x_2P denoted X_1 and X_2 respectively. This may be termed a first
 20 intra-entity key exchange.

Next, member A_1 computes $r = x_1P + x_2P$ and similarly, entity A_2 computes $r = x_2P + x_1P$. Thus, establishes an intra-entity shared key available and containing a contribution from each member of the entity.

The entity A transmits the intra-entity shared key r to the entity B with whom it wishes to
 25 establish a common key K .

Next, member A_1 computes a short term intra-entity public key s_1 using its short term private key and long term private key combined with a function f of the intra-entity public key, that is $s_1 = x_1 + a_1 f(r) \pmod n$, where f is typically a hash function such as SHA-1 and n is the order of the curve. Similarly, member A_2 computes its intra-entity public key $s_2 = x_2 + a_2 f(r)$
 30 $\pmod n$.

The entity A computes an entity or group short term public key, which is derived from a summation of the intra-entity public key of each member $s = s_1 + s_2 = x_1 + x_2 + (a_1 + a_2) f(r) \bmod (n)$.

The entity B similarly computes the analogous information using its own public and private keys using the same computations performed by entity A. Thus, each member of B computes a intra-entity shared key \bar{r} using the short term public keys of each of the members and \bar{r} is forwarded to entity A. Next, each of the members in B compute their own intra-entity public key $t_i = y_i + b_i f(\bar{r}) \bmod (n)$ and computes the group short-term public key $t = t_1 + t_2$.

The entity A then computes a value K which is the shared key between the entities A and B by retrieving the long term public key, bP , of entity B and computing $K = s(\bar{r} + (bP) f(\bar{r})) = s(t)P$. The entity B also retrieves the long term public key aP of entity A and computes K using t , r , and aP , i.e. $K = t(r + aP.f(r)) = t(s)P$.

Consequently, if a member of the entity A, either A_1 or A_2 , is not present in the scheme then the group short term public key, s , changes, as does the value for K. Therefore, communication with entity B would not be successful without establishing a new session. Similarly, if either B_1 or B_2 is not present in the scheme then the group short term public key, t , changes, altering the value of K. In this case, communication with A would not be successful without establishing a new session.

Accordingly, the present protocol ensures that all members of each entity must participate in the scheme and no sub-collection of either entity can impersonate its entire entity.

Although the above scheme has been described with respect to the elliptic curve systems which is an additive group, it may analogously be used in multiplicative groups. Furthermore the above protocol although exemplified with two members per entity, may be generalized where each party or entity consists of a collection of members say $A = \{A_1, A_2, \dots, A_n\}$ and $B = \{B_1, B_2, \dots, B_m\}$ where m is not necessarily equal to n and at least one of m or n is at least two (that is, not both A and B consist of one individual).

Although the invention has been described with reference to certain specific embodiments, various modifications thereof will be apparent to those skilled in the art without departing from the spirit and scope of the invention as outlined in the claims appended hereto.

THE EMBODIMENTS OF THE INVENTION IN WHICH AN EXCLUSIVE PROPERTY OR PRIVILEGE IS CLAIMED ARE DEFINED AS FOLLOWS:

1. A method for generating a shared secret value between entities (A,B) in a data communication system, one or more of said entities having a plurality of members (A_i, B_i) for participation in said communication system, each member having a long term private key and a corresponding long term public key said method comprising the steps of:
- (a) generating an entity long term private key and corresponding entity long term public key for each entity by combining the long term private and public keys of each members of the entity.
 - (b) generating a short term private and a corresponding short term public key for each of the members;
 - (c) exchanging short term public keys of the members within an entity;
 - (d) for each member:
 - i. computing an intra-entity shared key by mathematically combining said short term public keys of each said member;
 - ii. computing an intra-entity public key by mathematically combining its short-term private key, the long term private key and said intra-entity shared key;
 - (e) for each entity combining intra-entity public keys to derive a group short-term public key;
 - (f) each entity transmitting its intra-entity shared key and its group short term public key to said other entities; and
 - (g) each entity computing a common shared key K by combining its group short term public key, with the intra-entity shared key, and an entity long term public key received from the other entity.
2. A method as defined in claim 1, said long term public key being derived from a generator point P and respective ones of said long term private keys.
3. A method as defined in claim 2, said step (a) including each member selecting a random integer x_i and multiplying said point P by x_i to obtain x_iP , the short term public key.

4. A method as defined in claim 3, said intra-entity-shared key being computed by summing said short term public keys.
5. A method as defined in claim 4, said intra-entity public key s_i being derived by computing $s_i = x_i + a_i f(\sum x_i P)$, where f is a hash function.
6. A method as defined in claim 5, said group short term public key being derived by computing $\sum s_i$.
7. A method as defined in claim 1, said long term public keys being derived from a generator g and respective ones of said long term private keys.
8. A method as defined in claim 7, said step (a) including the step of each member selecting a random integer (x_{ij}) and exponentiating a function $h(g)$ including said generator to a power $g(x_{ij})$ to obtain the short term public key $X_{ij} = h(g)^{g(x_{ij})}$.
9. A method as defined in claim 8, said intra-entity shared key (X_i) being computed by each entity multiplying each of its short-term public keys X_{ij} together.
10. A method as defined in claim 1, including the step of exchanging the entity long term public key between entities.
11. A method as defined in claim 10, each entity computing a common shared key K by combining its group short term public key (S_i), with the intra-entity shared key (\bar{X}_i), and an entity long term public key received from the other entity.

ABSTRACT

This invention relates to a method for generating a shared secret value between entities in a data communication system, one or more of the entities having a plurality of members for participation in the communication system, each member having a long term private key and a corresponding long term public key. The method comprises the steps of generating a short term private and a corresponding short term public key for each of the members; exchanging short term public keys of the members within an entity. For each member then computing an intra-entity shared key by mathematically combining the short term public keys of each the members computing an intra-entity public key by mathematically combining its short-term private key, the long term private key and the intra-entity shared key. Next, each entity combines intra-entity public keys to derive a group short-term S_i public key; each entity transmitting its intra-entity shared key and its group short term public key to the other entities; and each entity computing a common shared key K by combining its group short term public key (S_i), with the intra-entity shared key (\bar{X}_i), and a group short term public (\bar{S}_i) key received from the other entities.

AD

$$\overbrace{A_1 \quad A_2}^A$$

$$\begin{array}{c} a_1 \quad a_2 \\ a_1 p \quad a_2 p \\ \downarrow \quad \downarrow \\ a p = a_1 p + a_2 p \end{array} \quad \xrightarrow{t b}$$

$$\begin{array}{c} x_1 \quad x_2 \\ x_1 p = x_1 \quad x_2 p = x_2 \\ \downarrow \quad \downarrow \\ r = x_1 + x_2 \end{array} \quad \xrightarrow{t b}$$

$$\begin{array}{c} s_1 = x_1 + a_1 f(r) \quad s_2 = x_2 + a_2 f(r) \\ \downarrow \quad \downarrow \\ s = s_1 + s_2 \end{array}$$

$$\text{compute } K = s(\bar{r} + (b p) f(\bar{r}))$$

FIG 2.

$$\overbrace{B_1 \quad B_2}^B$$

$$\begin{array}{c} b_1 \quad b_2 \\ b_1 p \quad b_2 p \\ \downarrow \quad \downarrow \\ b p = b_1 p + b_2 p \end{array} \quad \xrightarrow{t A}$$

$$\begin{array}{c} y_1 \quad y_2 \\ y_1 p = y_1 \quad y_2 p = y_2 \\ \downarrow \quad \downarrow \\ \bar{r} = y_1 + y_2 \end{array} \quad \xrightarrow{t A}$$

$$\begin{array}{c} t_1 = y_1 + b_1 f(\bar{r}) \quad t_2 = y_2 + b_2 f(\bar{r}) \\ \downarrow \quad \downarrow \\ t = t_1 + t_2 \end{array}$$

$$\text{compute } K = t(r + (b p) f(\bar{r}))$$